



## Frontiers of HPC: Unified Parallel C

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Some of this material is derived from notes written by Behzad Salami (M.Sc. U. of Guelph)



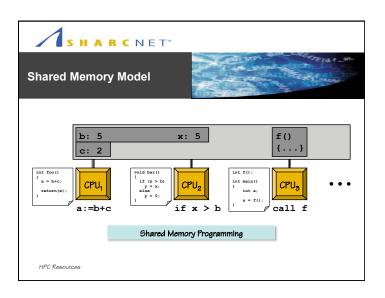


- Parallel programming models are not in short supply
  - methodology dictated by underlying hardware organization
    - shared memory systems (SMP)
    - · distributed memory systems (cluster)
  - trade-offs in ease of use; complexity
- · Unified Parallel C:
  - an open standard for a uniform programming model
  - distributed shared memory
  - fundamentals
  - a new set of trade-offs; is it worth it?

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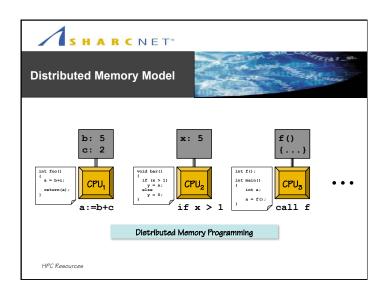
- Traditional parallel programming abstractly takes one of two forms, depending on how memory is able to be referenced by the program
- MIMD Models: Multiple Instruction Multiple Data
  - Shared memory
    - processors share a single physical memory
    - · programs can share blocks of memory between them
    - issues: exclusive access, race conditions, synchronization, scalability
  - Distributed memory
    - · unique memory associated with each processor
    - · issues: communication is explicit, communication overhead





- Shared memory programming benefits from the ability to handle communication implicitly
  - using the shared memory space
  - fundamentals of programming in SMP environments is relatively straightforward
    - issues typically revolve around exclusive access and race conditions
- · Common SMP programming paradigms:
  - POSIX threads (pthreads)
  - OpenMP

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```
#include <stdio.h>
#include *stdio.h>
#include "pthread.h"

void output (int *);

int main(int argc, char *argv[])

{
   int id;
   pthread_t thread[atoi(argv[1]);

   for (id = 0; id < atoi(argv[1]); id++)
      pthread_create(&thread[id], NULL, (void *)output, (void *)&id);

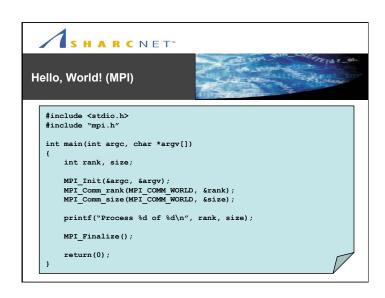
   return(0);
}

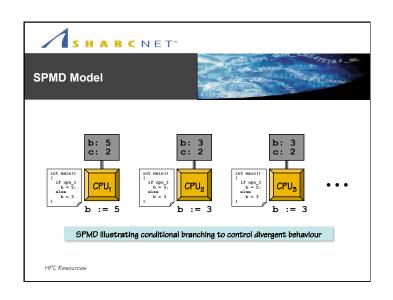
void output(int *thread_num)

{
   printf("Hello, world! from thread %d\n", *thread_num);
}
```



- · Communication is handled explicitly
  - processes send data to one another over an interconnection network
  - communication overhead limits granularity of parallelism
  - conforms to the strengths of traditional computing hardware so scalability can be excellent
- The past is littered with the corpses of distributed programming models
  - the modern standard: MPI







- Single Program Multiple Data (SPMD)
  - special case of MIMD model
  - many processors executing the *same* program
    - conditional branches used where specific behaviour required on the processors
  - shared/distributed memory organization
- · MPI and UPC explicitly use this model

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- C language extensions for HPC programming on large scale parallel systems
  - attempts to provide a single, uniform programming model for SMP and cluster-based machines
  - superset of C (any C program is automatically a valid UPC program)
- Explicitly parallel SPMD model
  - the same program runs on all processors
  - Global Address Space (GAS) language; an attempt to balance
    - convenience (threads)
    - performance (MPI, data layout)



#### What is UPC? (cont.)

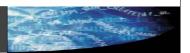


- Single shared address space
  - variables can be accessed by any process, however are physically associated with one
    - · hugely convenient from the programmer's perspective
    - what implications does this have?
  - most of UPCs complexity comes from the way it handles pointers and shared memory allocations
- Front-end/back-end organization allows for great flexibility in implementation
  - high speed interconnect, SMP, MPI, etc.
- · Well suited to parallelizing existing serial applications

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#### Major Features: Basics



- Multi-processing abstractly modeled as threads
  - pre-defined variables THREADS, MYTHREAD available at run-time
- new keyword: shared
  - defines variables available across all processes
  - affinity (physical location of data) can also be specified
    - scalars (affinity to process 0)
    - · cyclic (per element)
    - block-cyclic (user-defined block sizes)
    - blocked (run-time contiguous affinity for "even" distribution)
  - can make affinity inquiries at run-time

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#### Hello, World! (UPC)



 Note: even for the simplest of examples, the implicit availability of THREADS and MYTHREADS variable reduces code volume dramatically (over pthreads or MPI)

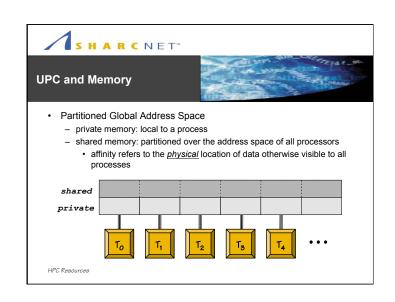
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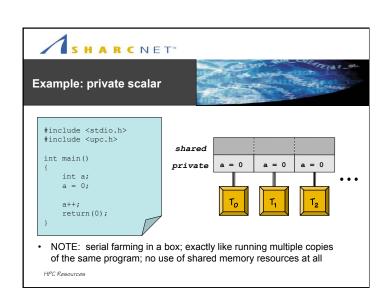


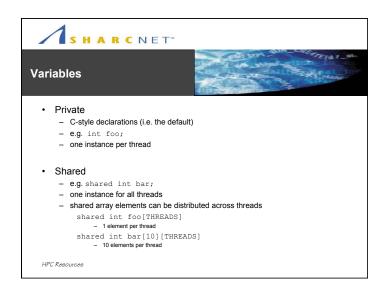
#### Major Features: Advanced

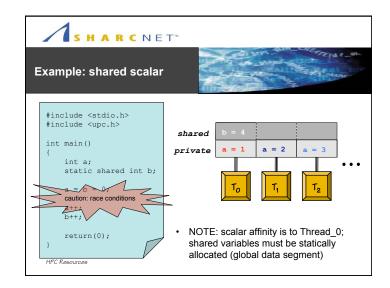


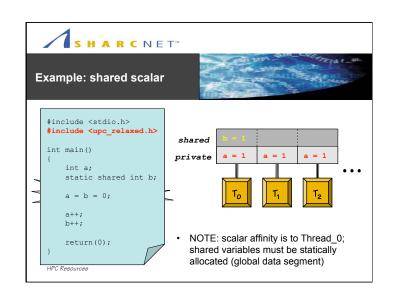
- Synchronization
  - locks (upc block t)
  - barrier (upc\_barrier, upc\_notify, upc\_wait)
  - memory fence (upc fence)
- · User-controlled consistency models
  - Per-variable, per-statement
  - strict
    - all strict references within the same thread appear in program order to all processes
  - relaxed
  - all references within the same thread appear in program order to the issuing process

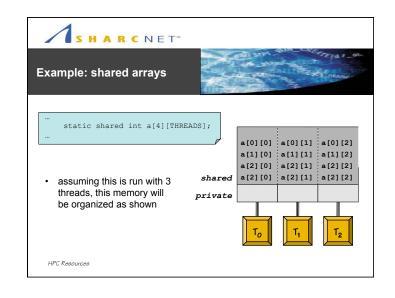


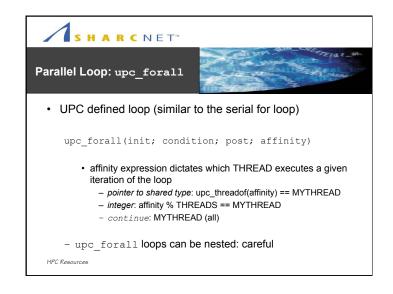


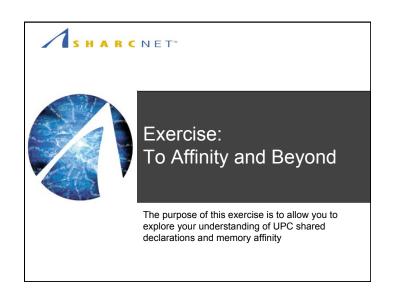








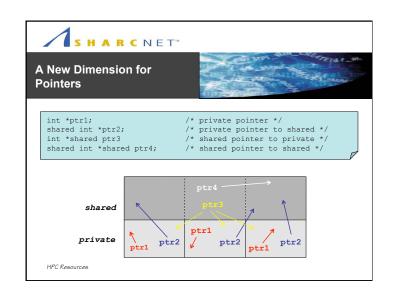






- The sharedarray.c file in ~dbm/public/ exercises/UPC implements a working demonstration of a shared allocation of a 2-D array of integers and the output of its contents by processor affinity (using the upc\_forall loop)
- Ensure that you understand the issue of processor affinity by changing the initialization of the array so that it too occurs in parallel, in the thread with affinity to that section of the array

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#### **Pointers**



· Declaration:

shared int \*a;

- a is a pointer to an integer that lives in the shared memory space
- we refer to the type of a as a pointer to shared

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#### Dynamic Memory Allocation



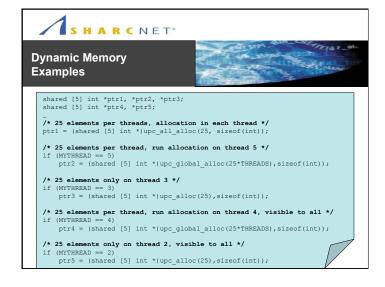
void \*upc\_all\_alloc(int num, int size);

- allocates num \* size bytes of data across all threads
- returns a pointer to the first element of the shared block of memory for all threads
- collective operations: called by all threads
- e.g. allocate 25 elements per thread



shared [5] int \*shared ptr2; /\* note use of block size \*/
...
ptr2 = (shared int \*)upc\_global\_alloc(25\*THREADS, sizeof(int));

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# Dynamic Memory Allocation (cont.) void \*upc\_alloc(int num, int size); - allocates num \* size bytes of data in the shared memory space, only on the calling thread (i.e. affinity issue) - returns a pointer to the first element of that shared block of memory for all threads - e.g. allocate 25 elements on Thread\_0, providing the reference to all threads shared [5] int \*shared ptr3; /\* note use of block size \*/

ptr3 = (shared int \*)upc alloc(25, sizeof(int));

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report results();

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Synchronization

upc\_barrier;

- all threads will wait until all threads execute the barrier

- very similar to MPl\_Barrier call (for those familiar)

- e.g.

if (MYTHREAD == 0)
 do\_initialization();

upc\_barrier;
/\* computation \*/
upc\_barrier;
if (MYTHREAD == 0)



#### Example/Exercise



- Let's consider a simple problem:
  - matrix multiplication
- · UPC refinement from serial code:
  - code it sequentially
  - data parallel: share/block data using THREADS
  - task parallel: organize parallel work by processor affinity using upc forall, THREADS and MYTHREAD
  - optimizing it by playing with memory affinity options

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# Exercise: The Matrix: Reparallelized

The purpose of this exercise is to illustrate the parallelization of a simple problem using UPC



#### **Exercise**



- The serialmatrix.c file in ~dbm/public/ exercises/UPC implements the classic matrix multiplication algorithm
- Parallelize this code using UPC by following the refinement steps outlined previously:
  - redeclare the 2D array to distribute it among threads
  - parallelize the work using affinity of data and upc forall
  - consider alternate memory affinity to optimize performance

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#### Summary

- Pro...
  - no need for explicit communication (on one hand: simpler code)
  - unified SMP/cluster model
  - compiler "optimized" communication
    - ensure optimal use of transport layer
    - relaxed consistency allows source reordering not possible in MPI
  - promise of scalability across paradigms
  - amenable to parallelizing existing serial code (for some domains)

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### • Con...

- no need for explicit communication (on the other: hidden complexity)
- introduces <u>new</u> complexity
  - distributed memory allocation
  - UPC advantages subverted by other language features (upc\_fence)
- language translation occurs at compile time limiting flexibility
- on some level: all the baggage of threads (race conditions, etc.) and MPI (collective communication, etc.)



#### **Food for Thought**



- Small study at U. of Guelph (B. Salami, 2005):
  - attempt to study "ease of programming" when using MPI/UPC, by those new to the programming model
  - not large enough to draw statistical conclusions, but some interesting observations
  - experienced undergraduate CS students:
    - MPI-trained group were challenged by synchronizing send/ recv calls, and issues of task parallelism
    - UPC-trained group were challenged by the pointer models and distributed memory allocation process; a more fundamental issue
    - · MPI-trained group got further with solutions
  - interesting implications for development in a research setting

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#### **Further Reading**



• Berkeley Unified Parallel C:

http://upc.lbl.gov/docs/

- examples
- reference documentation
- software (free)

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#### Discussion



- We haven't even seriously addressed issues of efficiency here
   what are your expectations?
- Distributed shared memory is a non-trivial problem
  - UPC is close to state of the art in terms of exposing a usable programming framework for DSM
    - significant complexity remains, some of it hidden from the programmer which can cause serious performance issues
- · Trading complexity
  - simple examples look promising; real code requires significant effort to do right; trade-offs in what requires the programmer's attention
- · Is UPC ready for prime-time?